# CS:APP Web Aside DATA:TMIN: Writing TMin in C* 

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June 5, 2012

## Notice

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## 1 The Situation

In Figure CS:APP2e-2.18 and in Problem CS:APP2e-2.21, we carefully wrote the value of TMin 32 as $-2147483647-1$. Why not simply write it as either -2147483648 or $0 \times 80000000$ ? Looking at the C header file limits.h, we see that they use a similar method as we have to write TMin $_{32}$ and $\operatorname{TMax}_{32}$ :

```
/* Minimum and maximum values a 'signed int' can hold. */
#define INT_MAX 2147483647
#define INT_MIN (-INT_MAX - 1)
```

Unfortunately, a curious interaction between the asymmetry of the two's complement representation and the conversion rules of C force us to write TMin $_{32}$ in this unusual way. Although understanding this issue requires us to delve into one of the murkier corners of the C language standards, it will help us appreciate some of the subtleties of integer data types and representations.
Consider the case of writing TMin $_{32}$ as -2147483648 and compiling the code on a 32 -bit machine, using the data sizes shown in Figure CS:APP2e-2.8. When the compiler encounters a number of the form $-X$, it first determines the data type and value for $X$ and then negates it. The value $2,147,483,648$ is too large to

[^0]| ISO C90 |  | ISO C99 |  |
| :--- | :--- | :--- | :--- |
| Decimal | Hexadecimal | Decimal | Hexadecimal |
| int <br> long <br> unsigned <br> unsigned long | int <br> unsigned <br> long <br> unsigned long | int <br> long <br> long long <br> unsigned | long <br> unsigned long <br> long long <br> unsigned long long |

Figure 1: Data types for representing integer constants. According to the language version and format (decimal or hexadecimal), the data type for a constant is given by the first type in the appropriate list that can represent the value.

| Word Size | ISO C90 |  | ISO C99 |  |
| :--- | :--- | :--- | :--- | :--- |
| Expression | -2147483648 | $0 \times 80000000$ | -2147483648 | $0 \times 80000000$ |
| 32 | unsigned | unsigned | long long | unsigned |
| 64 | long | unsigned | long | unsigned |

Figure 2: Data types resulting from constant expressions for TMin $_{32}$. According to the language version and format (decimal or hexadecimal), we can get three different data types for the two expressions, including cases where the value is positive.
represent as an int, since this value is one larger than $T M a x_{32}$ (the asymmetry strikes!). The compiler tries to determine a data type that can represent this value properly. It proceeds down one of the lists shown for the decimal cases in Figure 1, depending on the language version. For the case of ISO C90, it proceeds from int to long to unsigned, only then finding a data type that can represent the number 2,147,483,648. As we will see in CS:APP2e-2.3.3, values 2,147,483,648 and $-2,147,483,648$ have the same bit representations as 32 -bit numbers, and so the resulting constant has data type unsigned and value 2147483648 . For the case of ISO C99, the compiler proceeds from int to long to long long, finally finding a data type that can represent the number $2,147,483,648$. With 64 bits, we can uniquely represent both $2,147,483,648$ and $-2,147,483,648$, and so the resulting constant has data type long long and value -2147483648 .
For hexadecimal constant $0 \times 80000000$ on a 32 -bit machine, the compiler proceeds in a similar fashion, following one of the lists for the hexadecimal cases in Figure 1. For both language versions, it first compares the number to $T M a x_{32}$ ( $0 \times 7$ FFFFFFF) and, since it is larger, decides that the value cannot be represented as an int. It next compares the number to $U M a x_{32}(0 \times F F F F F F F F)$ and, since it is smaller, chooses an unsigned representation. It therefore yields a constant with data type unsigned and value $0 \times 80000000$ (or, equivalently, $2,147,483,648$ ).

Things work a bit differently on a 64-bit machine. For both language versions, the decimal form yields a constant with data type long (64-bits) and value $-2,147,483,648$, while the hexadecimal form yields a constant with type unsigned and value $0 \times 80000000$ (or, equivalently, 2,147,483,648).
All of these variations can be summarized by the table shown in Figure 2. For the cases where the result has type long or long long, the constant is negative, but it is 64 bits long. For the cases where the result
has type unsigned, the constant is positive and 32 bits long. These outcomes can be demonstrated by the following code:

```
int dcomp = (-2147483648<0);
int hcomp = (0x80000000< 0);
```

These lines of code attempt to express $T M i n_{32}$ as a decimal or hexadecimal constant and test whether the value is less than zero. Depending on the compiler version and word size, we find that the value of dcomp can be either 0 or 1 , indicating that the decimal constant can be either negative or positive, while the value of hcomp is consistently 0 , indicating that the hexadecimal constant is consistently positive. Our seemingly simple task of writing TMin $_{32}$ as a constant is more difficult than might be expected!

## Practice Problem 1:

Consider the following code:

```
int dtmin = -2147483648;
int dcomp2 = (dtmin < 0);
int htmin = 0x80000000;
int hcomp2 = (htmin < 0);
```

We compile this code on both 32 -bit and 64 -bit machines using two's complement representations of integers, and we try it for both language standards ISO-C90 and ISO-C99. In all cases, we consistently get value 1 for both dcomp2 and hcomp2, and further tests verify that dtmin and htmin indeed equal $\operatorname{TMin}_{32}$. Explain why this code does not have the compiler and language sensitivities we saw for the earlier code example.

## 2 Implications

For many programs, the ambiguities caused by different word sizes and language standards would not affect program behavior (for example, see Problem 1.) Nonetheless, we can now appreciate why the convention of writing TMin $_{32}$ as $-2147483647-1$ yields a more desirable result. Since 2147483647 is the value of $T M a x_{32}$, it can be represented as an int, and hence there is no need to invoke the conversion rules of Figure 1.

## Practice Problem 2:

Suppose we try to write $T_{M i n}^{32}$ as $-0 \times 7$ FFFFFFF-1. Would the C compiler generate a constant of type int for both 32 and 64 -bit machines and for both versions of the C language standard? Explain.

## Practice Problem 3:

You wish to write a succinct expression for $\operatorname{TMin}_{w}$, where $w$ is the number of bits in data type long int. Since the size of this data type varies from one machine to another (see Figures CS:APP2e-2.8 and CS:APP2e-2.9), you decide to make use of the sizeof operation, so that the expression will yield $T M i n_{w}$ as long as $w$ is a multiple of 8. You also use a trick, to be covered in Section CS:APP2e-2.3.6, that shifting a number left by 3 is the same as multiplying it by 8 .
Your first attempt at this code is:

```
/* WARNING: This code is buggy */
/* Shift 1 over by 8*sizeof(long) - 1 */
1L << sizeof(long)<<3 - 1
```

You test your code on a 32-bit machine, and find that the expression evaluates to 64 .
A. Explain why this happened.
B. What value would the expression yield on a 64 -bit machine?
C. Make minimal modifications to the expression so that it evaluates correctly.

## Solutions to Practice Problems

## Problem 1 Solution: [Pg. 3]

In making the assignment to integer variables dtmin and htmin, we implicitly cast the value to a 32 -bit, two's complement integer. This yields the value $-2,147,483,648$ regardless of whether or not the constant value is signed or unsigned, or whether it is 32 or 64 bits.

## Problem 2 Solution: [Pg. 3]

Yes, this would work as expected regardless of word size and language standard. Since $0 \times 7$ FFFFFFF is equal to $T M a x_{32}$, it will represent this value with data type int. The resulting expression therefore has data type int.

## Problem 3 Solution: [Pg. 3]

This is a classic example of failing to consider the operator precedence rules in C. As mentioned in Section CS:APP2e-2.1.10, addition and subtraction have higher precedence than shifting, and shifting associates to the left.
A. Consider the case where data type long requires 4 bytes. Then the expression is equivalent to $1 \ll 4 \ll 3-1$, which evaluates as $(1 \ll 4) \ll 2$, yielding 64 .
B. When long requires 8 bytes, we would have $1 \ll 8 \ll 3-1$ which evaluates as $(1 \ll 8) \ll 2$, yielding 1024.
C. The problem can be fixed with just one set of parentheses:

```
/* Shift 1 over by 8*sizeof(long) - 1 */
1L << (sizeof(long)<<3) - 1
```

We could also replace sizeof(long) <<3 by $8 *$ sizeof (long), and the higher precedence of multiplication would ensure correct expression evaluation. In fact, this would make the code more readable, and the resulting machine-level code would be identical.


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